# Learning-Augmented Online Algorithms 

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Motivating example: binary search
$n$ elements

| 8 | 11 | 14 | 16 | 18 | 25 | 30 | 36 | 40 | 43 | 46 | 49 | 50 | 53 | 54 | 56 | 59 | 60 | 63 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |

$$
q=16
$$

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## Introductio <br> Ski rental \& extensions <br> Motivating example: binary search

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Prediction: position $h(q) \quad$ Error: $\eta=\mid h(q)-$ index $(q) \mid$

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Motivating example: binary search


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Practical applications [KraskaBeutalChiDeanPolyzotis'18]

## Properties we seek



Algorithms are oblivious to $\eta$
Prediction $h$ should be learnable, e.g., compact

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## Classic" Beyond worst-case analysis

Future instance: $\quad X_{1} ; X_{2} ; X_{3} ; X_{4} ; X_{5} ; \ldots$

© Strong assumptions, needs some perfect information (oracle)

## HERE: no assumption on the predictor

allows plug-and-play predictors


$57.7 \%$ confidence


## Landscape preview

## learned <br> prediction ned

untrusted augmented
learning-aug
advice
Denomination
online conversion capacity scaling searchMTS
multiple-expert
query policies expert S ecretar Y iner tree primal-dual $S$ MQatching eneroy wi w reng ving energy ki or reng ing no follod best neighbor coaching on aching tasmana
bloom filters routing graph algorithms management

100s publications since 2018
https://algorithms-with-predictions.github.io/

## Most common framework used



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Objective: "minimize" competitive ratio $c_{A}(\eta)=\max _{I} \frac{\operatorname{cost}_{A}(I)}{\operatorname{OPT}(I)}$

$c_{A}(0)$


## Outline

## (1) Introduction

(2) Ski rental \& extensions

## (3) Preview of Paging and Graph Algorithms

## 4. Conclusion


cost to buy skis
daily rent price
? \# ski days (unknown)

Online algorithm: minimize $\max \frac{\operatorname{cost}}{\mathrm{OPT}}$

Best deterministic algo: buy at day $\approx b$

Worst-case $=$ stop after day $b$ :

- $\mathrm{Opt}=b$
- cost $=2 b$
- $\Longrightarrow$ competitive ratio $=2$


What should $h$ predict ?

- $)^{-} h \longrightarrow 0 / 1$ : rent or buy ? cannot measure $\eta$
- ;) $h \longrightarrow x$ with $\eta=|h-x|$

What should the algorithm do ?

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$\rightarrow$ - $h \longrightarrow 0 / 1$ : rent or buy ? cannot measure $\eta$
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What should the algorithm do ?
NAIVE: if $h \geq b$ then buy on day 1 else rent forever

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- ;) $h \longrightarrow x$ with $\eta=|h-x|$

What should the algorithm do ?
NAIVE: if $h \geq b$ then buy on day 1 else rent forever

## Lemma

The competitive ratio of NAIVE is $1+\eta$ / Opt.


Intuition: if X , we should not buy at day 1
How long should we rent? depends on the predictor's "trustworthiness"

$\operatorname{SkiPred}(\lambda):-\operatorname{If} h \geq b:$ rent $\lceil\lambda b\rceil$ days $\downarrow$ Else: rent $\lceil b / \lambda\rceil$ days

| rent | b | h |
| :--- | :--- | :--- | time

## Theorem

$\operatorname{SKIPRED}\left(\frac{1}{2}\right)$ is: $\min \left(3,1.5+2 \cdot \frac{\eta}{\mathrm{OPT}}\right)$-competitive.

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rent $\quad$ h b $\longrightarrow$ time

## Theorem

$\operatorname{SKIPRED}(\lambda)$ is: $\min \left(\frac{1+\lambda}{\lambda},(1+\lambda)+\frac{1}{1-\lambda} \cdot \frac{\eta}{\mathrm{OPT}}\right)$-competitive.


Classic randomized ski rental $\rightarrow \frac{e}{e-1} \approx 1.58$-competitive


Introduction Ski rental \& extensions

## Randomized ski rental

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## Theorem

There is a $O\left(\min \left(\frac{1}{1-e^{\lambda}}, \frac{\lambda}{1-e^{-\lambda}}\left(1+\frac{\eta}{\mathbf{O P T}^{\prime}}\right)\right)\right)$-competitive algorithm.

e.g., $\lambda=1 / 2$


ONLINE


Lower bounds:

- Randomized: matches UB [WeiZhang'20]
- Deterministic: LB a bit lower but


Dynamic Power Management: shift of focus

- focus first on 2 states
- il "free" over many rounds (experts framework)


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New tradeoff:


$$
\text { in cost }=\frac{1 \times 4}{n \mathrm{x}} \cdot \text { Opt }+\square \cdot \eta
$$

Use the whole prediction
Example for $\mathrm{a} \approx 1.16$-consistent, 0.38 -smooth solution:


## Multi-round ski rental [Antoniadis Coester Elias Polak Simon '21]


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(4) Conclusion

| $k=4 \quad$ misses: 1 | $\square$ |
| :---: | :---: |
| pages $\in\{A, B, \ldots, F\}$ | $\square$ |


| $k=4 \quad$ misses: 2 |  |
| :---: | :---: |
| pages $\in\{A, B, \ldots, F\}$ |  |
|  |  |
|  |  |

12
A B

| $k=4 \quad$ misses: 2 |  |
| :---: | :---: |
|  |  |
|  | pages $\in\{A, B, \ldots, F\}$ |
|  |  |
|  |  |

123
A B A

| $k=4 \quad$ misses: 3 |  |
| :---: | :---: |
|  |  |
|  | pages $\in\{A, B, \ldots, F\}$ |
|  |  |
|  |  |

$$
\begin{array}{cccc}
1 & 2 & 3 & 4 \\
\text { A } & \text { B } & \text { A } & \text { C }
\end{array}
$$

## Paging with predictions <br> [LykourisVassilvitski'18]

$$
\begin{array}{|c|c|}
k=4 \quad \text { misses: } 4 & \\
\hline \mathrm{C} \\
\hline & \mathrm{~B} \\
\hline & \mathrm{~A} \\
\hline
\end{array}
$$

$$
\begin{array}{ccccc}
1 & 2 & 3 & 4 & 5 \\
A & B & A & C & D
\end{array}
$$

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## Paging with predictions <br> [LykourisVassilvitski'18]

| $k=4 \quad$ misses: 6 | $D$ |
| :---: | :---: |
| pages $\in\{A, B, \ldots, F\}$ | E |
|  | A |
|  |  |

$$
\begin{array}{cccccccc}
1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \\
\mathrm{~A} & \mathrm{~B} & \mathrm{~A} & \mathrm{C} & \mathrm{D} & \mathrm{E} & \mathrm{~F} & \mathrm{~A}
\end{array}
$$

## Paging with predictions <br> [LykourisVassilvitski'18]

| $k=4 \quad$ misses: 7 | $D$ |
| :---: | :---: |
| pages $\in\{A, B, \ldots, F\}$ | E |
|  | A |
|  |  |

$$
\begin{array}{ccccccccc}
1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 \\
A & B & A & C & D & E & F & A & B
\end{array}
$$

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| $k=4 \quad$ misses: 7 | $D$ |
| :---: | :---: |
| pages $\in\{A, B, \ldots, F\}$ | E |
|  | A |
|  |  |

$$
\begin{array}{cccccccccc}
1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 \\
A & B & A & C & D & E & F & A & B & E
\end{array}
$$

## Paging with predictions <br> [LykourisVassilvitski'18]

| $k=4 \quad$ misses: 8 | $D$ |
| :---: | :---: |
| pages $\in\{A, B, \ldots, F\}$ | E |
|  | A |
|  |  |

$$
\begin{array}{ccccccccccc}
1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 & 11 \\
A & B & A & C & D & E & F & A & B & E & F
\end{array}
$$

## Paging with predictions <br> [LykourisVassilvitski'18]

| $k=4 \quad$ misses: 8 | $F$ |
| :---: | :---: |
| pages $\in\{A, B, \ldots, F\}$ | E |
|  | $A$ |
|  |  |

$$
\begin{array}{ccccccccccc}
1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 & 11 \\
A & B & A & C & D & E & F & A & B & E & F
\end{array}
$$



## Q: What to predict?

Lookahead (next q requests)

- © useless in the worst case

Strong Lookahead
(next requests until $q$ distinct)

- © huge, hard to predict

Next arrival time of the current request

- :) compact, enough to compute Opt, arguably learnable
- error $\eta_{i}$ at round $i$ : distance between predicted time and actual time combined error $\eta=\sum \eta_{i}$.
- $\Longrightarrow$ get $\mathrm{a} \approx \min \left(\log k, \log \frac{\eta}{\mathrm{OPT}}\right)$-competitive algorithm



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## Paging with predictions - Overview of models

Predictions $=$ time of next occurrence of current page

- Lykouris Vassilvitskii (2021 JACM); Rohatgi (SODA 2020); Wei (APPROX/RANDOM 2020)

Predictions $=$ all pages before next occurrence of current page

- Jiang Panigrahi Su (ICALP 2020)

Predictions $=$ state of ОРт (which pages in cache)

- Antoniadis Coester Elias Polak Simon (ICML 2020)

Multiple predictors - time of next occurrence of current page

- Emek Kutten Shi (ITCS 2020)

Prediction queries - obtain next occurrence of any page in cache

- Im Kumar Petety Purohit (ICML 2022)

Succinct predictions $=1$ bit of information ( $\approx$ to evicted or not)

- Antoniadis Boyar Eliáš Favrholdt Hoeksma Larsen Polak Simon (ICML 2023)


## A general error measure for graph algorithms

[Azar Panigrahi Touitou, Online Graph Algorithms with Predictions, SODA 2022]


Figure 1: An illustration of metric error with outliers. The figur
Error measure for predicting a set of points in a metric space $\eta=(D, \Delta)$ : $\mathrm{D}=$ transportation distance ;
$\Delta=\#$ outliers

## Theorem

For the Steiner Tree and Facility Location problems, if the error of the predicted input (resp. set of terminal and set of clients) is ( $D, \Delta$ ), there is an algorithm of cost at most $O(\log \Delta)$ Opt $+O(D)$.

## Introduction Ski rental \& extensions <br> Faster matching via learned duals

## [Dinitz Im Lavastida Moseley Vassilvitskii NeurIPS 2021] <br> [Chen Silwal Vakilian Zhang ICML 2022]

## Theorem

Given a weighted bipartite graph and predicted dual $\hat{y}$, there exists an algorithm that finds a minimum weight perfect matching in time $O\left(m \sqrt{n}+(m+n \log n)\left\|y^{*}-\hat{y}\right\|_{0}\right)$, where $y^{*}$ is an optimal dual solution.

Main ideas:

- use predicted dual as a warm start for the Hungarian algorithm
- if this dual is not feasible, adapt it
- actually design the algorithm that get the solution faster if the prediction error is small
- show that the duals can be learned from samples of a probabilistic input


## Irtroduction

## (2) Ski rental \& extensions

(3) Preview of Paging and Graph Algorithms

4 Conclusion

## Conclusion

## Take-back messages

- 



- fresh algorithm concepts
- relevant link ML - algorithms


## Newer questions

- improve running time
- ensure learnability (e.g., PAC-learnability) / $\eta \approx$ loss function
- extensive experiments including ML predictors
- multiple predictors
- ll wrt renowned heuristics ?

